Blockchain: Magic, Mechanics and Methods

Stephen J. Mildenhall November 28, 2018



Blockchain: Magic and Marketing

Your Business Problems

Your Business Problems

Customer Experience

- Confusing products
- High expenses

Your Business Problems

Customer Experience Confusing products • High expenses Your Problem 1 Your Problem 2 Your Problem 3

Your Business Problems

Universal Solution!

Customer Experience
Confusing products
High expenses
Your Problem 1
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•
Your Problem 2
•
•
Your Problem 3
•

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Customer Experience

- Confusing products
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Your Problem 1

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Your Problem 2

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Your Problem 3

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Universal Solution!

Blockchain delivers...

- Best customer experience
- Immutable record
- Enables collaboration
- One view of truth
- End duplicate reconciliation
- Lower costs
- Eliminates fraud
- Regulatory compliance
- Product innovation
- Quick to market
- ·_____
 - · _____

Your Business Problems

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- Confusing products
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Your Problem 3

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Any sufficiently advanced technology is indistinguishable from magic.

Arthur C. Clarke

Definition

Blockchains are **distributed** digital **ledgers** of **cryptographically signed transactions** that are grouped into **blocks**. Each block is **cryptographically linked** to the previous one after **validation** and undergoing a **consensus decision**, making it **tamper evident**. As new blocks are added, older blocks become more **difficult to modify**. New blocks are **replicated** across copies of the ledger within the network, and any **conflicts** are **resolved automatically** using established rules.

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Description

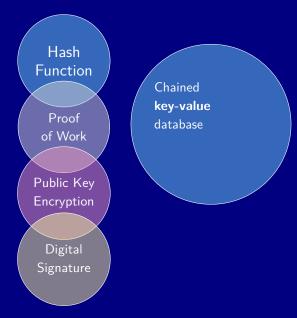
Components

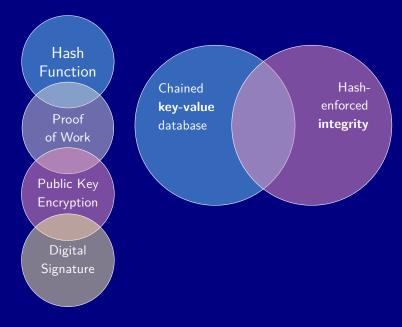
- Distributed database
- Ledger
- Cryptographically...
- ...Signed transactions
- ...Linked (chained)
- Consensus Validation

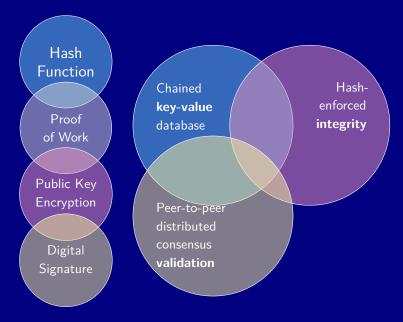
Characteristics

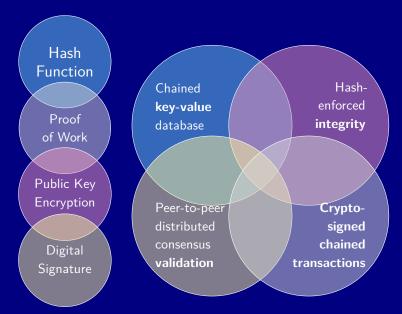
- No authority
- High availability
- Replicated, robust
- Tamper evident
- Difficult to modify
- Conflicts resolved







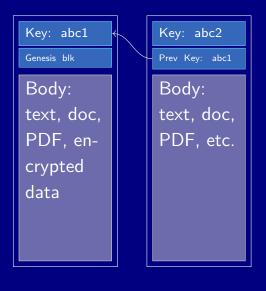


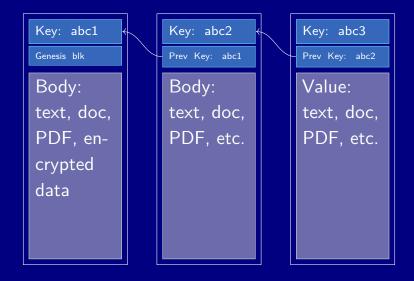


Key: abc1 Body: text, doc, PDF, encrypted data

Key: abc1 Body: text, doc, PDF, encrypted data

Key: abc2 Body: text, doc, PDF, etc.





Ingredient: Hash Functions

A hash H maps data of arbitrary size to a fixed size such that

- H(x) is an easy to compute, deterministic function
- If $x \neq y$ then $H(x) \neq H(y)$ with high probability
- H(x) appears random over its range as x varies
- IT hash function: first five letters of last name + first letter first name
- J. Smith problem
- Phone, zip, social, . . .

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Cryptographic Hash Function

- Given y it is **very hard** to find x with H(x) = y
- Fuggedaboutit hard

SHA256 Cryptographic Hash Function

import hashlib

hashlib.sha256(b'The quick brown fox jumps over the lazy dog').hexdigest() >>> 'd7a8fbb307d7809469ca9abcb0082e4f8d5651e46d3cdb762d02d0bf37c9e592'

hashlib.sha256(b'The quick brown fox jumps over the lazy dog.').hexdigest() >>> 'ef537f25c895bfa782526529a9b63d97aa631564d5d789c2b765448c8635fb6c'

- Output = **very** large integer, between 0 and $2^{256} \approx 10^{77}$
- Specify input and output formats very carefully
- Probability of J. Smith collision: not even a Dumb and Dumber chance

The Birthday Problem and Hash Collisions

- Birthday problem: 23 people for 50/50 chance of same birthday
- Number of documents before p probability of collision given a hash space size of N is $\approx \sqrt{2Np}$ for small p^1
- For SHA256, $N = 2^{256} = 10^{77}$ is very large
- $\, \blacksquare \,$ A 10^{-3} collision probability requires about 1.5×10^{37} documents, enough for
 - Every person on earth to...
 - Compute 1 billion hashes per second...
 - For five times the age of the universe

¹E.g. for birthday problem p=1/2, N=365 and $\sqrt{2Np}=19$. Approximation relies on $p\approx -\log(1-p)$, only true for smaller p. Using $(-2N\log(1-p))^{1/2}=22.49$ is very close to correct answer, 23.

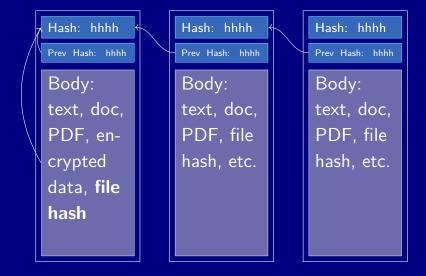
Ingredient 2: Hash-Enforced Integrity



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Hash: hhhh Hash: hhhh Prev Hash: hhhh Prev Hash: hhhh Body: Body: text, doc, text, doc, PDF, en-PDF, file crypted hash, etc. data, file hash

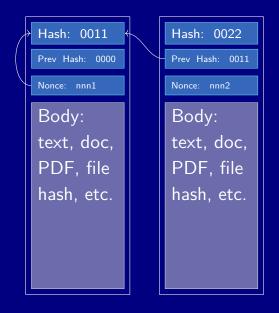
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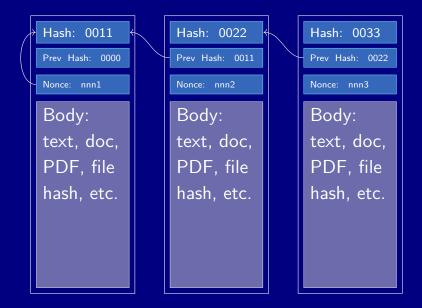
Ingredient 3: Distributed Validation and Proof-Of-Work



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Proof of Work and Bitcoin Mining = Compute Hashes

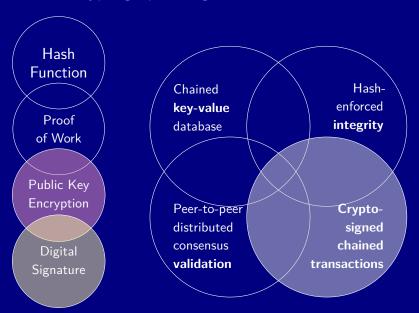
```
import hashlib
running min = 2**256
ans = \Pi
base = b'The quick brown fox jumps over the lazy dog'
for nonce in range(2000000000): # 2 billion
    h = hashlib.sha256(nonce.to_bytes(4, byteorder='big') + base).hexdigest()
    n = int(h, 16)
    if n < running_min:</pre>
        running_min = n
        ans.append([nonce, n])
        print(f'{nonce:12,d} {n:077}')
```

Proof of Work and Bitcoin Mining = Compute Hashes

Nonce Hash 29115579230639891023898657467946481563928575965694753738500728003067276450760 21833633896494697913657452817095065461049276120751755746016193921330837964982 17391853960576662285627567225372501697536440120814058733709287576654299269058 00 491741673371171570027367996335736784622791320015893772572199978008540614786 817 00 207113148484537618144604663416437589440289273319027116671254033065643419132 827 000 35029650895291714754047120679492927968250654901817817434081241936987361735 3.292 000 30590294895123458493702891527069975442971551875566805022772671084264919745 6.362 000 23157006908555232018903879877754051315219896322661305099606253143774488785 7,634 000 11843095073522994422561274720857316931066719486382550615573171404879921966 22.034 0000 6045160764465103256154815045992679930360222615550766779824452388654984639 32,737 0000 3218718010716516807246023638919032202673987969434384430166215105132280583 43.078 0000 3066940367111277087798394765784480513227788830972580117541505418890948712 50.740 00000 344804005194498392473362848134761831134304453202875173759130216105619080 260,109 00000 149043122808237032345561872905133216060467384369910593113997965062602336 610,827 000000 25441204939268765420155917698735840343496809686969451042687651132777655 3,553,698 000000 12372585984995238023081534031026808791454761919139475665549030259593011 16,603,005 0000000 4682308792444739613119316155033986067282587356863979013510780284611482 45.767.445 0000000 4295135810439807939037487563409966578108755229939605598485594694500274 56.389.936 0000000 1219890553970511010693160459086914039690075265862677724048817741406404 186,599,009 00000000 741733398915175814111679160159562329641666849535152212310255158283708 187,060,155 00000000 129027976973068678554136418237268320708790839626316760173444080235551 209,437,773 00000000 46418792192972977622708878642780226280538977482131916077098153688658 554,751,705 00000000 38492057003517052607600918969310106371482316138230835578404460555913 1,724,412,865 000000000 20951411954830677538112338658105096359813168232452740277675602777590

- Current network hash rate 4×10^{19} hashes per second
- Electricity consumption = Austria
- Block hash: 0x 0000 0000 0000 0001 a841 86ab c5df

Dissect: Cryptographic Ingredients



Discrete Logarithm Problem

Discrete logarithm problem says

given
$$g^a \equiv n \pmod{p}$$
 can't find a

is a one-way function

mod p means remainder after dividing by prime p

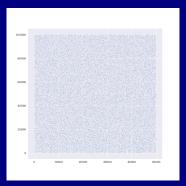
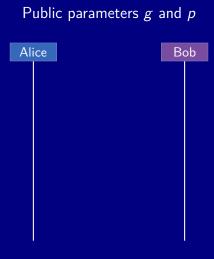
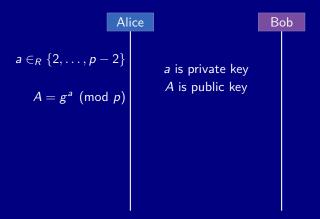


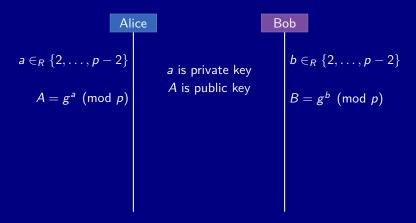
Figure 1: Powers of 3 modulo 100043; $100042 = 2 \times 50021$ is twice a prime.



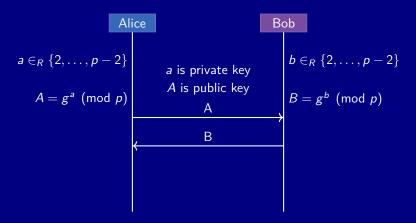
Public parameters g and p



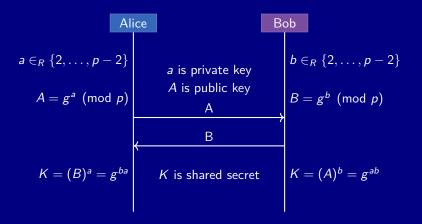
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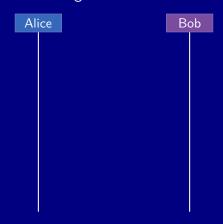
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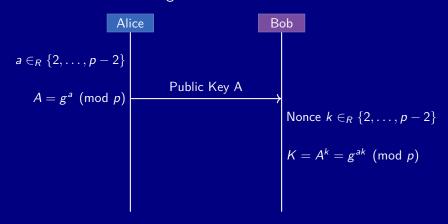
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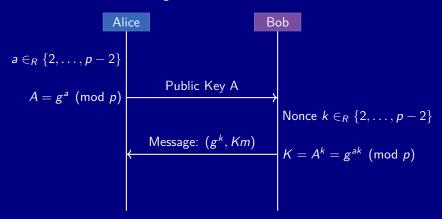
Public parameters g and pSend message m from Bob to Alice



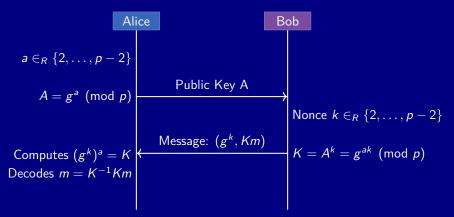
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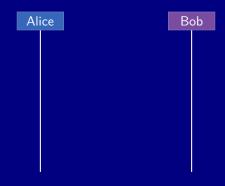


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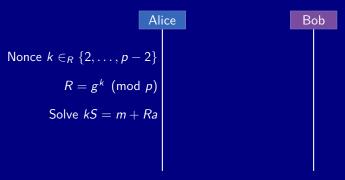


 g^k conveys information about k but shields its value; K hides message m

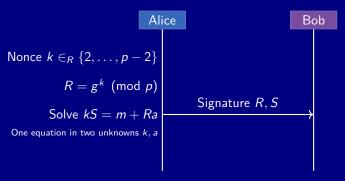
Alice to sign message m, Bob to verify $g, p, A = g^a, m$ all public, a is secret



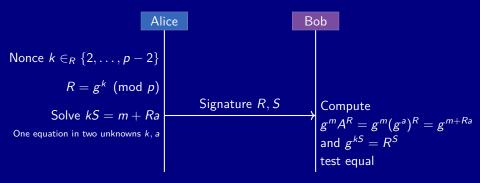
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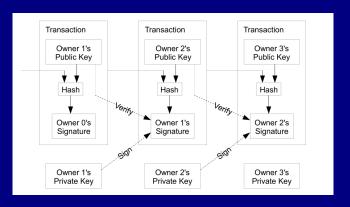
If Alice does not know a she can't find R, S to solve $R^S = g^m A^R$

Powerful Properties of Digital Signature

- Signer authentication: verifier assured that signature has been created only by sender who possess the corresponding secret private key
- Message integrity: if message modified, signature fails; signature tamper evident
- Non-repudiation: existence of signature proves it came from sender; sender cannot repudiate signing in future
- Wet ink signatures can be forged; document can be altered; signature can be denied

Ingredient 4: Double-spend mechanism

- Bitcoin ledger tracks coin ownership
- Owners can endorse to new owners in cryptographically secure manner
- Public pseudonymous chain of ownership

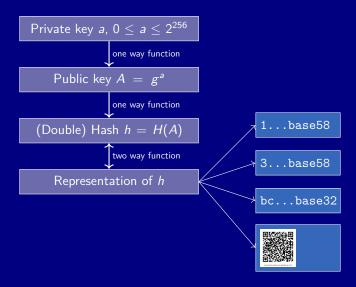


What is a Bitcoin Public Address?

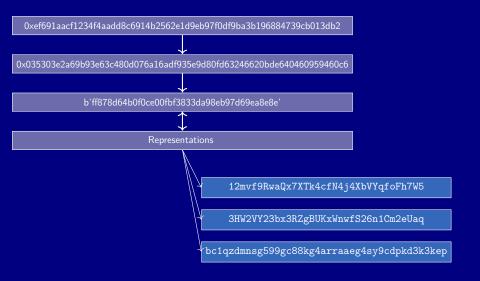


Figure 2: Donations gratefully received.

What is a Bitcoin Public Address?



What is a Bitcoin Public Address?



If You Know What You Are Doing...

Load into Bitcoin Core Client and get addresses via WIF compressed representation of private key importprivkey L5F6PZo9h2RJnGGvztwWEUnwYH1eWhpv63Z5qQEZgqxcy364nBCQj yourName getaddressesbyaccount yourName "12mvf9RwaQx7XTk4cfN4j4XbVYqfoFh7W5", "3HW2VY23bx3RZgBUKxWnwfS26n1Cm2eUaq", "bc1qzdmnsg599gc88kg4arraaeg4sy9cdpkd3k3kep"

Discovery: Solution in Search of a Problem

Using ingredients...

- Hash functions
- Public/private keys
- Digital signatures
- Chained blocks
- Chained transactions
- A clever incentive reinforcing recipe

We have created a...

- Distributed...
- Available...
- Public/unsuppressable...
- Immutable database
- No central authority
- Trust between strangers
- Digital scarcity

Must discover applications requiring new features...

not just trust = legal contract...

not just highly available = DNS, GAFA...

we've built a tank of the database world...





You Could Drop the Kids Off at School in a Tank



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You Could Drop the Kids Off at School in a Tank



...and you'd likely end up with a...



...SQL database

Capabilities and Refinements Are In Conflict

Between	and	there is a conflict
Obvious TTP	Blockchain	Trusted third party administers SQL DB
Public	Permissioned	Coordinate without blockchain
Open source	Governance	Uncoordinated open network = forks
Privacy	Verifiability	Information needed to verify transactions
Trust	Performance	Low/no trust = poor performance
Access	Efficiency	Guaranteed access, distributed = expensive
PII	Public	Expectation of privacy
PII	Immutable	GDPR Right to be forgotten
Me	Everyone else	Coordination or technology problem?

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 Confidential transactions can keep the amount and type of assets transferred visible only to participants in the transaction (and those they choose to reveal the blinding key to), while still cryptographically guaranteeing that no more coins can be spent than are available



Self-sovereign identity: now that it's possible, it's inevitable.

Humanity deserves digital identity that is permanent, portable, private and completely secure; in other words: self-sovereign.

Shortcomings in the internet's original design made this impossible, at a cost of <u>trillions</u> each year. Today, the invention of distributed ledger technology makes self-sovereign digital identity a possibility for the first time.

Now that self-sovereign identity is possible, it's inevitable. And it's going to change everything.

Identity is the Killer App

Self-Sovereign Identity and Decentralized Identifiers (DIDs)

- Permanent
- Resolvable
- Cryptographically Verifiable
- Decentralized
- Verifiable credentials
- Store data on edge devices, no central stores of PII
- User control of own data
- GDPR

Drummond Reed, Decentralized Identifiers (DIDs) The Fundamental Building Block of Self-Sovereign Identity https://goo.gl/Au4uBx

Zero Knowledge Proofs

- It is possible to verify information without revealing it: using a zero knowledge proof
- Read-only access, read-act-forget
- Where's Waldo with a mat
- Alibaba's cave

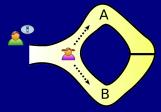


Figure 3: Alibaba's cave example: it is possible to prove you know something without revealing it.

Industry Consortia and Alliances

- R3: distributed ledger, banking; created Corda
- RiskBlock Alliance (The Institutes)
- B3i: blockchain Insurance Industry Initiative (London)
- AAIS: openIDL = open Insurance Data Link, regulatory data reporting
- Alastria national blockchain system

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- Cooperation and coordination is a social problem not a technology problem!

Commercial Solutions

- Etherisc: travel and other insurances on Ethereum
- Everledger: registry for diamonds and other real assets, an identity solution
- NodalBlock: customer on-boarding, document commitment

Deployment Options

- Ethereum network smart contract, "world computer"
- Bitcoin network, did:btcr, Lightning network
- Hyperledger
- Corda open-source DLT/blockchain platform
- Private forks of open source solutions

Conclusions

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Blockchain Pros

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- Enables unimagined solutions
- Perfect for identity problems

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- Amazing technical capability the Internet circa 1995
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Blockchain Cons

- Slow, expensive database
- Cyber/real-world interface about ambiguity not smart contracts
- Coordination still required