## Blockchain: Magic, Mechanics and Methods

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## Blockchain: Magic and Marketing

[Your Industry] Blockchain Opportunity Slide
[Your Industry] Blockchain Opportunity Slide
Your Business Problems
[Your Industry] Blockchain Opportunity Slide
Your Business Problems
Customer Experience

- Confusing products
- High expenses
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Your Problem 1

Your Problem 2

Your Problem 3
[Your Industry] Blockchain Opportunity Slide
Your Business Problems
Universal Solution!

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## Universal Solution!

## Blockchain delivers...

- Best customer experience
- Immutable record
- Enables collaboration
- One view of truth
- End duplicate reconciliation
- Lower costs
- Eliminates fraud
- Regulatory compliance
- Product innovation
- Quick to market
$\bullet$
- $\qquad$
[Your Industry] Blockchain Opportunity Slide

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Any sufficiently advanced technology is indistinguishable from magic.

Arthur C. Clarke

## Definition

Blockchains are distributed digital ledgers of cryptographically signed transactions that are grouped into blocks. Each block is cryptographically linked to the previous one after validation and undergoing a consensus decision, making it tamper evident. As new blocks are added, older blocks become more difficult to modify. New blocks are replicated across copies of the ledger within the network, and any conflicts are resolved automatically using established rules.

Blockchain Technology Overview, Yaga et al (2018), NIST

## Definition

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## Description

## Components

- Distributed database
- Ledger
- Cryptographically...
- ...Signed transactions
- ...Linked (chained)
- Consensus Validation


## Characteristics

- No authority
- High availability
- Replicated, robust
- Tamper evident
- Difficult to modify
- Conflicts resolved

Dissect: Magical Ingredients \& Recipe


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Ingredient 1: Chained Key-Value (Distributed) Database

```
Key: abc1
```

Body:
text, doc, PDF, encrypted data

Ingredient 1: Chained Key-Value (Distributed) Database


Key: abc2

Body: text, doc, PDF, etc.

Ingredient 1: Chained Key-Value (Distributed) Database


Ingredient 1: Chained Key-Value (Distributed) Database


## Ingredient: Hash Functions

A hash $H$ maps data of arbitrary size to a fixed size such that

- $H(x)$ is an easy to compute, deterministic function
- If $x \neq y$ then $H(x) \neq H(y)$ with high probability
- $H(x)$ appears random over its range as $x$ varies
- IT hash function: first five letters of last name + first letter first name
- J. Smith problem
- Phone, zip, social, ...


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Cryptographic Hash Function

- Given $y$ it is very hard to find $x$ with $H(x)=y$
- Fuggedaboutit hard


## SHA256 Cryptographic Hash Function

import hashlib
hashlib.sha256(b'The quick brown fox jumps over the lazy dog').hexdigest() >>> 'd7a8fbb307d7809469ca9abcb0082e4f8d5651e46d3cdb762d02d0bf37c9e592'
hashlib.sha256(b'The quick brown fox jumps over the lazy dog.').hexdigest() >>> 'ef537f25c895bfa782526529a9b63d97aa631564d5d789c2b765448c8635fb6c'

- Output $=$ very large integer, between 0 and $2^{256} \approx 10^{77}$
- Specify input and output formats very carefully
- Probability of J. Smith collision: not even a Dumb and Dumber chance


## The Birthday Problem and Hash Collisions

- Birthday problem: 23 people for 50/50 chance of same birthday
- Number of documents before $p$ probability of collision given a hash space size of $N$ is $\approx \sqrt{2 N p}$ for small $p^{1}$
- For SHA256, $N=2^{256}=10^{77}$ is very large
- A $10^{-3}$ collision probability requires about $1.5 \times 10^{37}$ documents, enough for
- Every person on earth to...
- Compute 1 billion hashes per second. . .
- For five times the age of the universe

[^0]Ingredient 2: Hash-Enforced Integrity


Ingredient 2: Hash-Enforced Integrity


Ingredient 2: Hash-Enforced Integrity


Ingredient 3: Distributed Validation and Proof-Of-Work

```
Hash: 0011
```

Prev Hash: 0000

Nonce: nnn1
Body:
text, doc,
PDF, file
hash, etc.

Ingredient 3: Distributed Validation and Proof-Of-Work


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## Proof of Work and Bitcoin Mining = Compute Hashes

import hashlib

```
running_min = 2**256
ans = []
base = b'The quick brown fox jumps over the lazy dog'
for nonce in range(2000000000): # 2 billion
    h = hashlib.sha256(nonce.to_bytes(4, byteorder='big') + base).hexdigest()
    n = int(h, 16)
    if n < running_min:
        running_min = n
        ans.append([nonce, n])
        print(f'{nonce:12,d} {n:077}')
```


## Proof of Work and Bitcoin Mining = Compute Hashes

Nonce

Hash
29115579230639891023898657467946481563928575965694753738500728003067276450760 21833633896494697913657452817095065461049276120751755746016193921330837964982 17391853960576662285627567225372501697536440120814058733709287576654299269058 00491741673371171570027367996335736784622791320015893772572199978008540614786 00207113148484537618144604663416437589440289273319027116671254033065643419132 00035029650895291714754047120679492927968250654901817817434081241936987361735 00030590294895123458493702891527069975442971551875566805022772671084264919745 00023157006908555232018903879877754051315219896322661305099606253143774488785 00011843095073522994422561274720857316931066719486382550615573171404879921966 00006045160764465103256154815045992679930360222615550766779824452388654984639 $0000 \quad 3218718010716516807246023638919032202673987969434384430166215105132280583$ 00003066940367111277087798394765784480513227788830972580117541505418890948712 00000344804005194498392473362848134761831134304453202875173759130216105619080 00000149043122808237032345561872905133216060467384369910593113997965062602336 00000025441204939268765420155917698735840343496809686969451042687651132777655 00000012372585984995238023081534031026808791454761919139475665549030259593011 00000004682308792444739613119316155033986067282587356863979013510780284611482 00000004295135810439807939037487563409966578108755229939605598485594694500274 00000001219890553970511010693160459086914039690075265862677724048817741406404 00000000741733398915175814111679160159562329641666849535152212310255158283708 00000000129027976973068678554136418237268320708790839626316760173444080235551 00000000046418792192972977622708878642780226280538977482131916077098153688658 00000000038492057003517052607600918969310106371482316138230835578404460555913 00000000020951411954830677538112338658105096359813168232452740277675602777590

- Current network hash rate $4 \times 10^{19}$ hashes per second
- Electricity consumption = Austria
- Block hash: 0x 00000000000000000051 a841 86ab c5df

Dissect: Cryptographic Ingredients


## Discrete Logarithm Problem

- Discrete logarithm problem says

$$
\text { given } g^{a} \equiv n(\bmod p) \text { can't find a }
$$

is a one-way function

- mod $p$ means remainder after dividing by prime $p$


Figure 1: Powers of 3 modulo $100043 ; 100042=2 \times 50021$ is twice a prime.

## Creating a Shared Secret

Public parameters $g$ and $p$


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Public/private pair $(A, a)$ are cryptographically linked but $a$ is hidden

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## ElGamel Public Key Encryption

Public parameters $g$ and $p$
Send message $m$ from Bob to Alice


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A=g^{a}(\bmod p) \underbrace{\text { Alice }}_{\text {Message: }\left(g^{k}, K m\right)} \underset{\text { Nublic Key } A}{ }
$$

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Send message $m$ from Bob to Alice


## Digital Signature

Alice to sign message $m$, Bob to verify $g, p, A=g^{a}, m$ all public, $a$ is secret


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| Alice |  | Bob |
| :---: | :---: | :---: |
| Nonce $k \in_{R}\{2, \ldots, p-2\}$ |  |  |
| $R=g^{k}(\bmod p)$ | Signature $R, S$ | Compute |
| One equation in two unknowns $k$, a |  | $\begin{aligned} & g^{m} A^{R}=g^{m}\left(g^{a}\right)^{R}=g^{m+R a} \\ & \text { and } g^{k S}=R^{S} \\ & \text { test equal } \end{aligned}$ |

If Alice does not know a she can't find $R, S$ to solve $R^{S}=g^{m} A^{R}$

## Powerful Properties of Digital Signature

- Signer authentication: verifier assured that signature has been created only by sender who possess the corresponding secret private key
- Message integrity: if message modified, signature fails; signature tamper evident
- Non-repudiation: existence of signature proves it came from sender; sender cannot repudiate signing in future
- Wet ink signatures can be forged; document can be altered; signature can be denied


## Ingredient 4: Double-spend mechanism

- Bitcoin ledger tracks coin ownership
- Owners can endorse to new owners in cryptographically secure manner
- Public pseudonymous chain of ownership



## What is a Bitcoin Public Address?



Figure 2: Donations gratefully received.

## What is a Bitcoin Public Address?



## What is a Bitcoin Public Address?



## If You Know What You Are Doing. . .

Load into Bitcoin Core Client and get addresses via WIF compressed representation of private key
importprivkey
L5F6PZo9h2RJnGGvztwWEUnwYH1eWhpv63Z5qQEZgqxcy364nBCQj yourName
getaddressesbyaccount yourName
[
"12mvf9RwaQx7XTk4cfN4j4XbVYqfoFh7W5", "3HW2VY23bx3RZgBUKxWnwfS26n1Cm2eUaq",
"bc1qzdmnsg599gc88kg4arraaeg4sy9cdpkd3k3kep"
]

Discovery: Solution in Search of a Problem

## Using ingredients... <br> - Hash functions <br> - Public/private keys <br> - Digital signatures <br> - Chained blocks <br> - Chained transactions <br> - A clever incentive <br> reinforcing recipe

We have created a...

- Distributed...
- Available...
- Public/unsuppressable...
- Immutable database
- No central authority
- Trust between strangers
- Digital scarcity

Must discover applications requiring new features... not just trust = legal contract... not just highly available = DNS, GAFA... we've built a tank of the database world...


## You Could Drop the Kids Off at School in a Tank



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Pros
Coolest kids in school

- Good if you run into trouble
- Don't need a road
- Park where ever you like


## You Could Drop the Kids Off at School in a Tank

Pros
Coolest kids in school

- Good if you run into trouble | -2 Cruising speed 30 mph |
| :--- |
| 2020 mph in 7 seconds |
- Park where ever you like


## Cons

Cost new $\$ 4.3$ million

- Don't need a road


## You Could Drop the Kids Off at School in a Tank

Pros

- Coolest kids in school
- Good if you run into trouble
- Don't need a road
- Park where ever you like


## Cons

Cost new $\$ 4.3$ million

- Cruising speed 30 mph
- 0 to 20 mph in 7 seconds
- Fuel economy 0.6 mpg
...and you'd likely end up with a...

...SQL database


## Capabilities and Refinements Are In Conflict

| Between | and | there is a conflict |
| :--- | :--- | :--- |
| Obvious TTP | Blockchain | Trusted third party administers SQL DB |
| Public | Permissioned | Coordinate without blockchain |
| Open source | Governance | Uncoordinated open network $=$ forks |
| Privacy | Verifiability | Information needed to verify transactions |
| Trust | Performance | Low/no trust = poor performance |
| Access | Efficiency | Guaranteed access, distributed = expensive |
| PII | Public | Expectation of privacy |
| PII | Immutable | GDPR Right to be forgotten |
| Me | Everyone else | Coordination or technology problem? |

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- Confidential transactions can keep the amount and type of assets transferred visible only to participants in the transaction (and those they choose to reveal the blinding key to), while still cryptographically guaranteeing that no more coins can be spent than are available


## Self-sovereign identity: now that it's possible, it's inevitable.

Humanity deserves digital identity that is permanent, portable, private and completely secure; in other words: self-sovereign.

Shortcomings in the internet's original design made this impossible, at a cost of trillions each year. Today, the invention of distributed ledger technology makes self-sovereign digital identity a possibility for the first time.

Now that self-sovereign identity is possible, it's inevitable. And it's going to change everything.

## Identity is the Killer App

## Self-Sovereign Identity and Decentralized Identifiers (DIDs)

- Permanent
- Resolvable
- Cryptographically Verifiable
- Decentralized
- Verifiable credentials
- Store data on edge devices, no central stores of PII
- User control of own data
- GDPR


## Zero Knowledge Proofs

- It is possible to verify information without revealing it: using a zero knowledge proof
- Read-only access, read-act-forget
- Where's Waldo with a mat
- Alibaba's cave


Figure 3: Alibaba's cave example: it is possible to prove you know something without revealing it.

## InsureTech Blockchain Applications

Industry Consortia and Alliances

- R3: distributed ledger, banking; created Corda
- RiskBlock Alliance (The Institutes)
- B3i: blockchain Insurance Industry Initiative (London)
- AAIS: openIDL = open Insurance Data Link, regulatory data reporting
- Alastria national blockchain system


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- Alastria national blockchain system
- Cooperation and coordination is a social problem not a technology problem!


## InsureTech Blockchain Applications

## Commercial Solutions

- Etherisc: travel and other insurances on Ethereum
- Everledger: registry for diamonds and other real assets, an identity solution
- NodalBlock: customer on-boarding, document commitment


## InsureTech Blockchain Applications

## Deployment Options

- Ethereum network smart contract, "world computer"
- Bitcoin network, did:btcr, Lightning network
- Hyperledger
- Corda open-source DLT/blockchain platform
- Private forks of open source solutions

Conclusions

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Blockchain Pros

- Amazing technical capabilitythe Internet circa 1995
- Enables unimagined solutions
- Perfect for identity problems


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- Amazing technical capabilitythe Internet circa 1995
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## Blockchain Cons

- Slow, expensive database
- Cyber/real-world interface about ambiguity not smart contracts
- Coordination still required


[^0]:    ${ }^{1}$ E.g. for birthday problem $p=1 / 2, N=365$ and $\sqrt{2 N p}=19$. Approximation relies on $p \approx-\log (1-p)$, only true for smaller $p$. Using $(-2 N \log (1-p))^{1 / 2}=22.49$ is very close to correct answer, 23.

